IDSDS. The Iterated Deletion of Strictly Dominated Strategies
$A$ is not a dominant straresy for Player 1
$B=$

Player 1


For Player $1 C$ is strictly dominated by $B$
After deleting $C$, now $D$ becomes strictly dominated by $F$ After deleting $D$, now $A$

$$
\begin{array}{ll}
R_{1}=\text { player } 1 & \text { is rational } \\
R_{2}=1,2 & B_{1} x=\text { player } 1 \text { believes } x \\
B_{2} x=1 " 2
\end{array}
$$

IDSDS. The Iterated Deletion of Strictly Dominated Strategies


$$
\begin{array}{cl}
R_{1} & B_{1} R_{2} \\
R_{2}, B_{2} R_{1} & B_{2} B_{1} B_{2} R_{1} \\
B_{1} B_{2} R_{1} & B_{2} B_{1} R_{2}
\end{array}
$$

At any stage delete strategies, that are either strictly of weakly dominated
IDWDS. The Iterated Deletion of Weakly Dominated Strategies

(2) because $R$ weakly dom. by $L$

IDWDS. The Iterated Deletion of Weakly Dominated Strategies

(2) bed. $L$ weanly dow. by $R$

Definition: at every stagehidentify all the strobegis, that are weakly or strictly dominated. Then delete cell of them ar the same IDWDS. The Iterated Deletion of Weakly Dominated Strategies time.

Player 2

|  | $L$ |  | L |  |
| :---: | :---: | :---: | :---: | :---: |
|  | 4 | 4 | 0 | 0 |

Repeat

For $1 \quad B$ is weakly dom. by $T$

IDSDS leave, the game unchanged
Nash equilibrium

$$
\begin{aligned}
& x \in S_{1}(=\text { set of strategies of Player 1) } \\
& y \in S_{2} \quad(\cdots \quad 2)
\end{aligned}
$$



Taus Nash equilubria:

$$
\begin{aligned}
& (A, E) \\
& (C, D)
\end{aligned}
$$

$(x, y)$ is a Nash equilibrium if
$\pi_{1}(x, y) \geq \pi_{1}(z, y)$ for every $z \in S_{1}$
$\Pi_{2}(x, y) \geq \Pi_{2}(x, w)$ for every $w \in S_{2}$
Player 2
C
D
Player 2


Player 3 chooses F
Player 3 chooses G
$(A, D, F)$ is a Nash equilibrium

| Player | $A$ | Player |  |  |  | 2 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 1 | 0 | 2 | 3 | 3 | 1 |
| 1 | B | 3 | 3 | 1 | 5 | 4 | 4 |
|  | C |  | 2 | 0 | 1 | 3 | 0 |

Player 2
C D

| Player | A | 2 | 2 | 2 | 4 | 3 | 6 |
| :---: | :---: | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{1}$ | B | 5 | 3 | 2 | 3 | 4 | 2 |

Player 3 chooses F 둔

Player 2
C D

Player 3 chooses G

Large game.
150 students in a class, they simultaneously ask for a grade (A, B or C); if $20 \%$ or less (ie. $\leq 30$ ) ask for an A then all requests are granted, otherwise they all get a C.

Selfish players
first set: exactly 30 choose $A, 120$ choose $B$
Second set: at least 32 choose A

Example with uncertain outcomes. A simple auction. There are two players, Charlie and Doreen. There is an object (e.g. a painting) which Charlie values at $\$ 120$ and Doreen values at $\$ 180$. Each player has to submit a bid of either $\$ 50$ or $\$ 80$. The highest bidder gets the object and pays his/her bid (the loser does not pay anything). If the bids are equal, a fair coin is tossed.

Outcomes: $\quad a$ Charlie wins and pays $\$ 50$
$b$ Charlie wins and pays $\$ 80$
c Doreen wins and pays $\$ 50$
d Doreen wins and pays $\$ 80$
Player's utility $=$ value - price paid (if wins, otherwise 0 )

Doreen (value: \$180)


Outcomes: $\quad a$ Charlie wins and pays $\$ 50$
$b$ Charlie wins and pays $\$ 80$
c Doreen wins and pays $\$ 50$
d Doreen wins and pays $\$ 80$
Player's utility $=$ value - price paid (if wins, otherwise 0 )

Doreen (value: \$180)

|  | bid \$50 | bid \$80 |
| :---: | :---: | :---: |
| Charlie bid $\$ 50$ | $\left(\begin{array}{ll}b & d \\ \frac{1}{2} & \frac{1}{2}\end{array}\right)$ | $d$ |
| (value: \$120) bid \$80 | $b$ | $\left(\begin{array}{ll}b & d \\ \frac{1}{2} & \frac{1}{2}\end{array}\right)$ |

Doreen

|  |  | bid $\$ 50$ |  |
| :---: | :---: | :---: | :---: |
| Charlie | bid $\$ 50$ | 35,65 | 0,100 |
|  | bid $\$ 80$ | 40,0 | 20,50 |
|  |  |  |  |

